Routing Basics

ISP Workshops



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Acknowledgements

- This material originated from the Cisco ISP/IXP Workshop Programme developed by Philip Smith & Barry Greene
- Use of these materials is encouraged as long as the source is fully acknowledged and this notice remains in place
- Bug fixes and improvements are welcomed
 - Please email workshop (at) bgp4all.com

Philip Smith

Routing Concepts

- Routers
- Routing
- Forwarding
- Some definitions
- Policy options
- Routing Protocols

What is a Router?

■ A router is a layer 3 device



- Used for interconnecting networks at layer 3
- A router generally has at least two interfaces
 - With VLANs a router can have only one interface (known as "router on a stick")
- A router looks at the destination address in the IP packet, and decides how to forward it

The Routing Table

- Each router/host has a routing table, indicating the path or the next hop for a given destination host or a network
- The router/host tries to match the destination address of a packet against entries in the routing table
- If there is a match, the router forwards it to the corresponding gateway router or directly to the destination host
- Default route is taken if no other entry matches the destination address

The Routing Table

Destination	Next-Hop	Interface
10.40.0.0/16	192.248.40.60	Ethernet0
192.248.0.140/30	Directly connected	Serial1
192.248.40.0/26	Directly connected	Ethernet0
192.248.0.0/17	192.248.0.141	Serial1
203.94.73.202/32	192.248.40.3	Ethernet0
203.115.6.132/30	Directly connected	Serial0
Default	203.115.6.133	Serial0

Typical routing table on a simple edge router

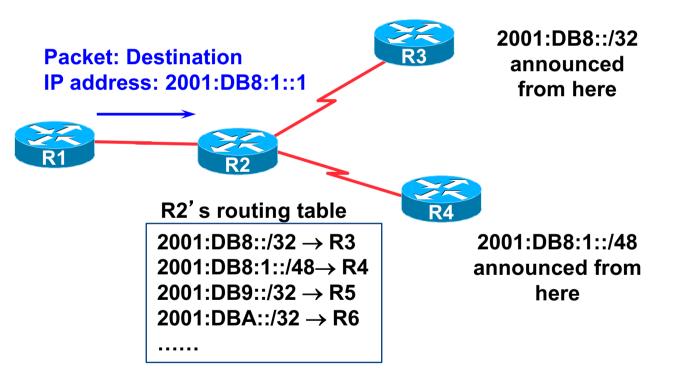
IP Routing – finding the path

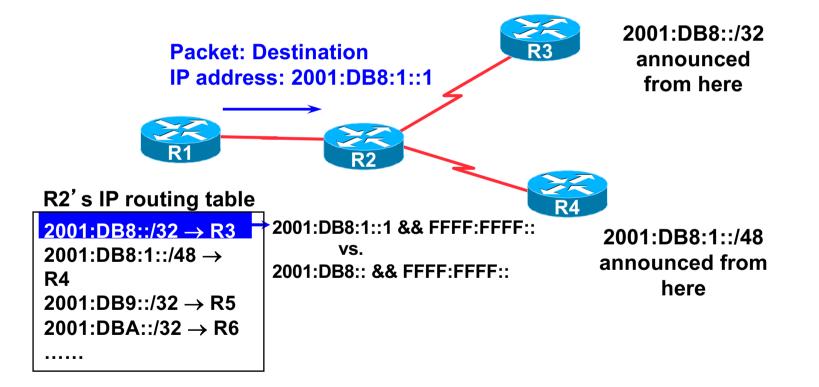
- Routing table entry (the path) is created by the administrator (static) or received from a routing protocol (dynamic)
- More than one routing protocol may run on a router
 - Each routing protocol builds its own routing table (Local RIB)
- Several alternative paths may exist
 - Best path selected for the router's Global routing table (RIB)
- Decisions are updated periodically or as topology changes (event driven)
- Decisions are based on:
 - Topology, policies and metrics (hop count, filtering, delay, bandwidth, etc.)

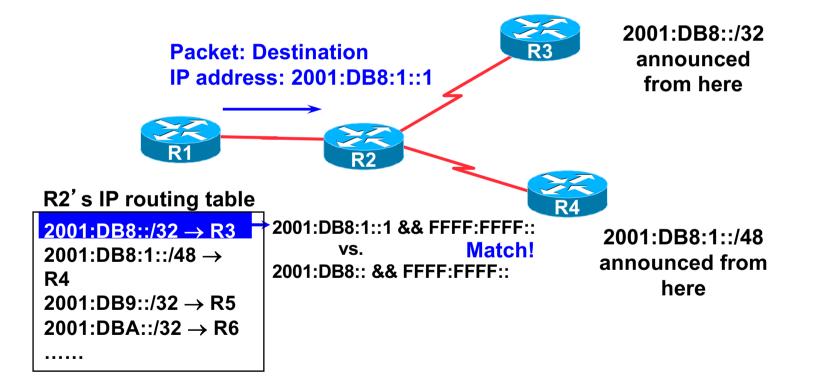
IP route lookup

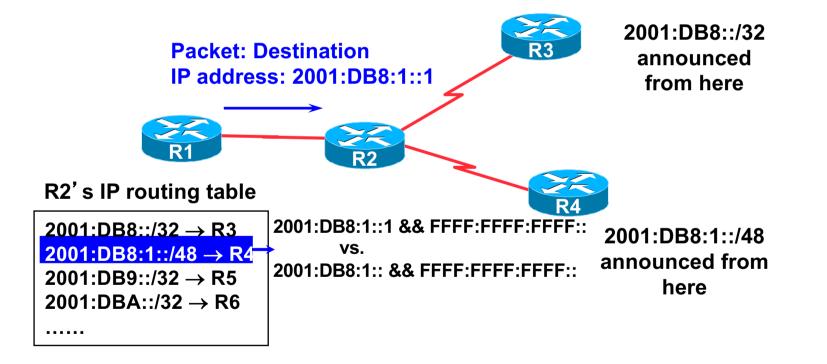
- Based on destination IP address
- "Iongest match" routing
 - More specific prefix preferred over less specific prefix
 - Example: packet with destination of 2001:DB8:1::1/128 is sent to the router announcing 2001:DB8:1::/48 rather than the router announcing 2001:DB8::/32.

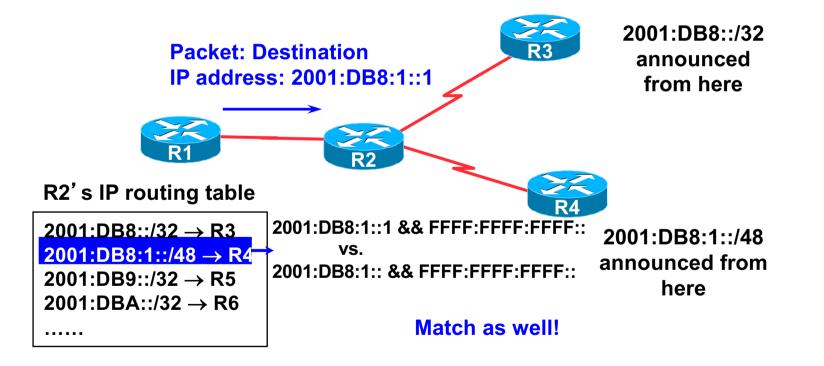
IP route lookup

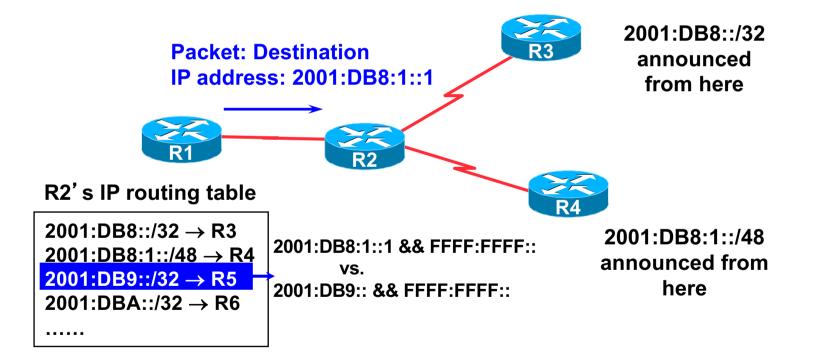


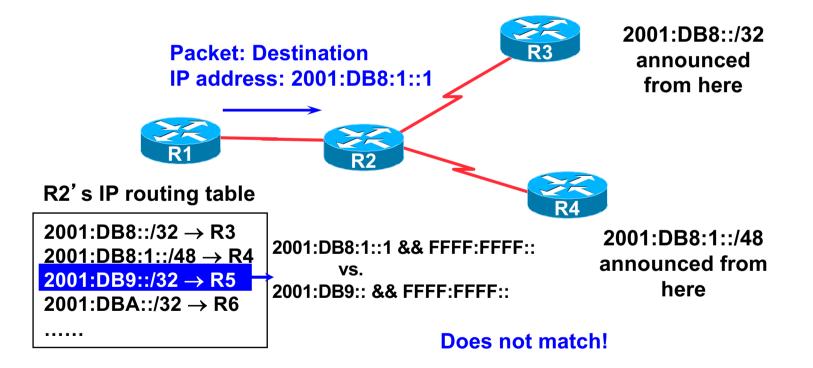


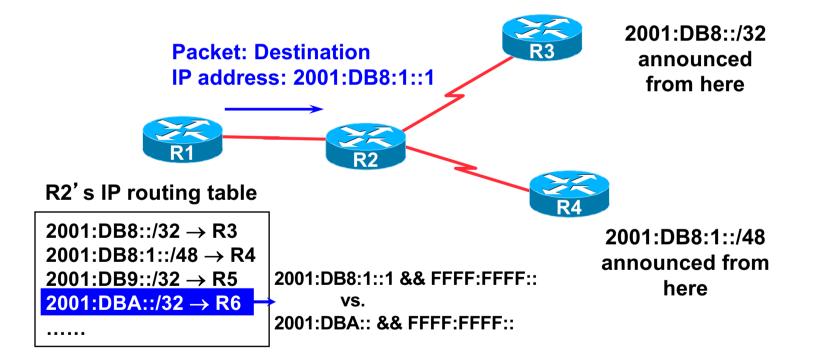


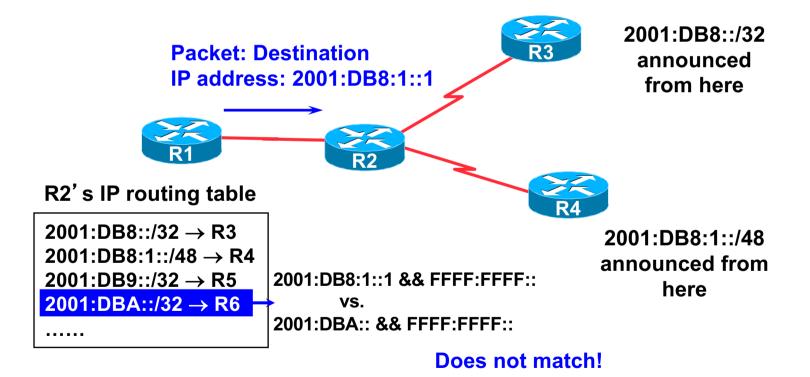


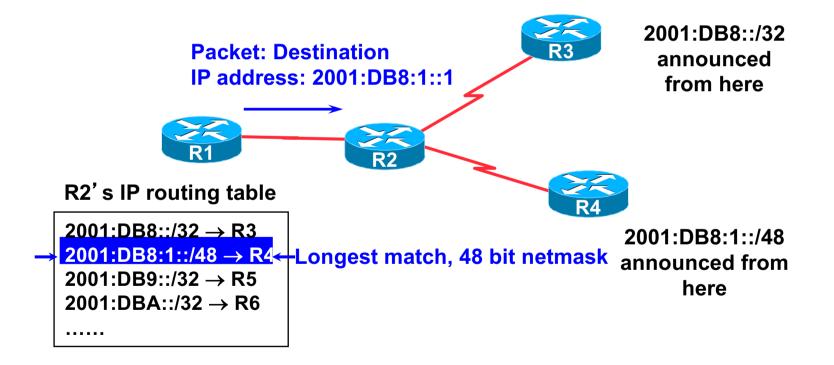


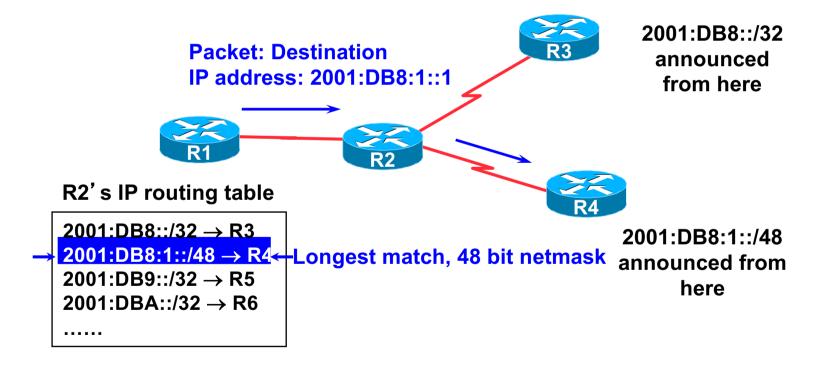






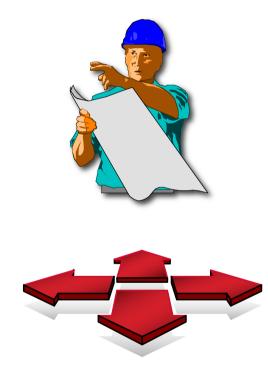






Routing versus Forwarding

- Routing = building maps and giving directions
- Forwarding = moving packets between interfaces according to the "directions"



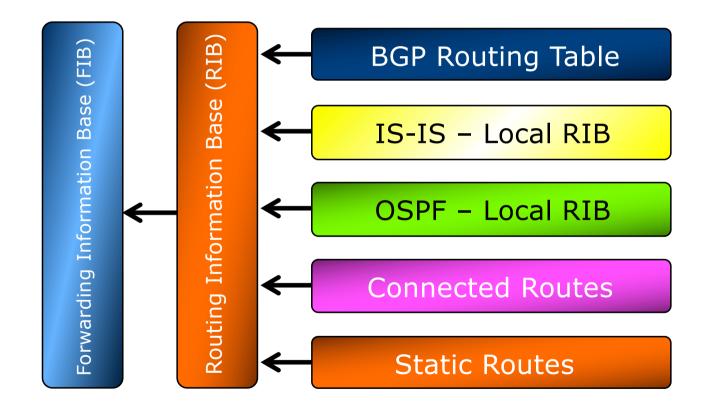
IP Forwarding

- Router decides which interface a packet is sent to
- Forwarding table populated by routing process

Forwarding decisions:

- destination address
- class of service (fair queuing, precedence, others)
- Iocal requirements (packet filtering)
- Forwarding is usually aided by special hardware

Routing Tables Feed the Forwarding Table



The FIB

□ FIB is the Forwarding Table

- It contains destinations, the interfaces and the next-hops to get to those destinations
- It is built from the router's Global RIB
- Used by the router to figure out where to send the packet
- Cisco IOS: "show ip cef"

The Global RIB

□ The Global RIB is the Routing Table

- Built from the routing tables/RIBs of the routing protocols and static routes on the router
 - Routing protocol priority varies per vendor see addendum
- It contains all the known destinations and the next-hops used to get to those destinations
- One destination can have lots of possible next-hops only the best next-hop goes into the Global RIB
- The Global RIB is used to build the FIB
- Cisco IOS: "show ip route"

Explicit versus Default Routing

- Default:
 - Simple, cheap (CPU, memory, bandwidth)
 - No overhead
 - Low granularity (metric games)
- Explicit: (default free zone)
 - Complex, expensive (CPU, memory, bandwidth)
 - High overhead
 - High granularity (every destination known)
- Hybrid:
 - Minimise overhead
 - Provide useful granularity
 - Requires some filtering knowledge

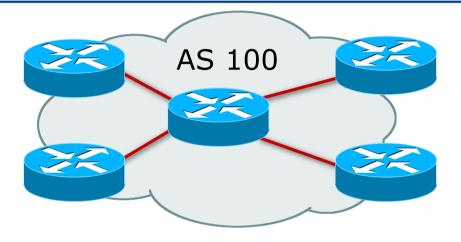
Egress Traffic

- How packets leave your network
- Egress traffic depends on:
 - Route availability (what others send you)
 - Route acceptance (what you accept from others)
 - Policy and tuning (what you do with routes from others)
 - Peering and transit agreements

Ingress Traffic

- How packets get to your network and your customers' networks
- Ingress traffic depends on:
 - What information you send and to whom
 - Based on your addressing and AS's
 - Based on others' policy (what they accept from you and what they do with it)

Autonomous System (AS)



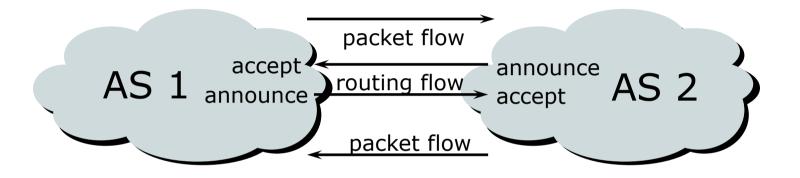
- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control

Definition of terms

Neighbours

- AS's which directly exchange routing information
- Routers which exchange routing information
- Announce
 - send routing information to a neighbour
- Accept
 - receive and use routing information sent by a neighbour
- Originate
 - insert routing information into external announcements (usually as a result of the IGP)
- Peers
 - routers in neighbouring AS's or within one AS which exchange routing and policy information

Routing flow and packet flow



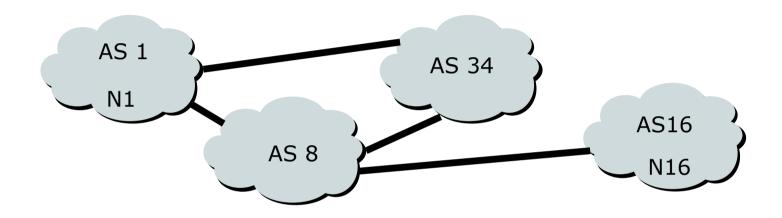
For networks in AS1 and AS2 to communicate:

AS1 must announce to AS2 AS2 must accept from AS1 AS2 must announce to AS1 AS1 must accept from AS2

Routing flow and Traffic flow

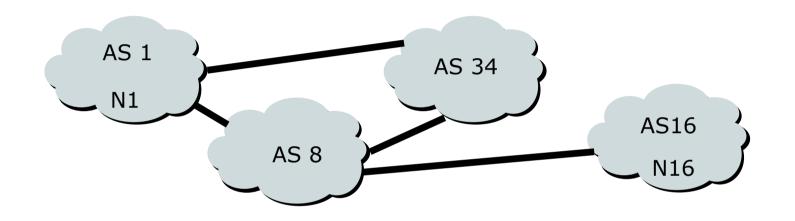
- Traffic flow is always in the opposite direction of the flow of Routing information
 - Filtering outgoing routing information inhibits traffic flow inbound
 - Filtering inbound routing information inhibits traffic flow outbound

Routing Flow/Packet Flow: With multiple ASes



- For net N1 in AS1 to send traffic to net N16 in AS16:
 - AS16 must originate and announce N16 to AS8.
 - AS8 must accept N16 from AS16.
 - AS8 must announce N16 to AS1 or AS34.
 - AS1 must accept N16 from AS8 or AS34.
- For two-way packet flow, similar policies must exist for N1

Routing Flow/Packet Flow: With multiple ASes

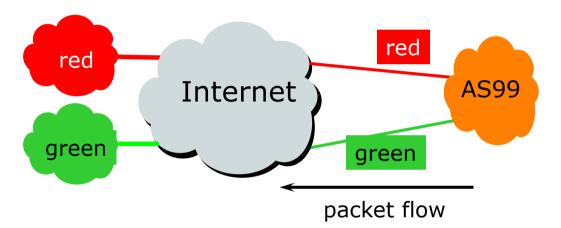


As more and more paths are implemented between sites it is easy to see how policies can become quite complex.

Routing Policy

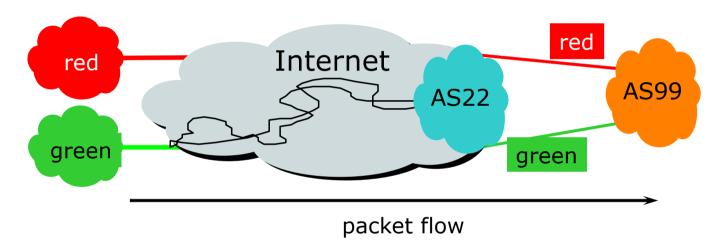
- Used to control traffic flow in and out of an ISP network
- ISP makes decisions on what routing information to accept and discard from its neighbours
 - Individual routes
 - Routes originated by specific ASes
 - Routes traversing specific ASes
 - Routes belonging to other groupings
 Groupings which you define as you see fit

Routing Policy Limitations



- AS99 uses red link for traffic to the red AS and the green link for remaining traffic
- To implement this policy, AS99 has to:
 - Accept routes originating from the red AS on the red link
 - Accept all other routes on the green link

Routing Policy Limitations



- AS99 would like packets coming from the green AS to use the green link.
- But unless AS22 cooperates in pushing traffic from the green AS down the green link, there is very little that AS99 can do to achieve this aim

Routing Policy Issues

December 2018:

- 58000 IPv6 prefixes & 728000 IPv4 prefixes
 Not realistic to set policy on all of them individually
- 62600 origin AS's
 - Too many to try and create individual policies for
- Routes tied to a specific AS or path may be unstable regardless of connectivity

Solution: Groups of AS's are a natural abstraction for filtering purposes

Routing Protocols

We now know what routing means... ...but what do the routers get up to? And why are we doing this anyway?

- Internet is made up of the ISPs who connect to each other's networks
- How does an ISP in Kenya tell an ISP in Japan what customers they have?
- And how does that ISP send data packets to the customers of the ISP in Japan, and get responses back
 - After all, as on a local ethernet, two way packet flow is needed for communication between two devices

- ISP in Kenya could buy a direct connection to the ISP in Japan
 - But this doesn't scale thousands of ISPs, would need thousands of connections, and cost would be astronomical
- Instead, ISP in Kenya tells his neighbouring ISPs what customers he has
 - And the neighbouring ISPs pass this information on to their neighbours, and so on
 - This process repeats until the information reaches the ISP in Japan

□ This process is called "Routing"

- The mechanisms used are called "Routing Protocols"
- Routing and Routing Protocols ensures that
 - The Internet can scale
 - Thousands of ISPs can provide connectivity to each other
 - We have the Internet we see today

- ISP in Kenya doesn't actually tell its neighbouring ISPs the names of the customers
 - (network equipment does not understand names)
- Instead, it has received an IP address block as a member of the Regional Internet Registry serving Kenya
 - Its customers have received address space from this address block as part of their "Internet service"
 - And it announces this address block to its neighbouring ISPs this is called announcing a "route"

Routing Protocols

- Routers use "routing protocols" to exchange routing information with each other
 - IGP is used to refer to the process running on routers inside an ISP's network
 - EGP is used to refer to the process running between routers bordering directly connected ISP networks

What Is an IGP?

- Interior Gateway Protocol
- Within an Autonomous System
- Carries information about internal infrastructure prefixes
- □ Two widely used IGPs:
 - OSPF
 - IS-IS

Why Do We Need an IGP?

ISP backbone scaling

- Hierarchy
- Limiting scope of failure
- Only used for ISP's infrastructure addresses, not customers or anything else
- Design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

What Is an EGP?

- Exterior Gateway Protocol
- Used to convey routing information between Autonomous Systems
- De-coupled from the IGP
- Current EGP is BGP

Why Do We Need an EGP?

- Scaling to large network
 - Hierarchy
 - Limit scope of failure
- Define Administrative Boundary
- Policy
 - Control reachability of prefixes
 - Merge separate organisations
 - Connect multiple IGPs

Interior versus Exterior Routing Protocols

Interior

- Automatic neighbour discovery
- Generally trust your IGP routers
- Prefixes go to all IGP routers
- Binds routers in one AS together

Exterior

- Specifically configured peers
- Connecting with outside networks
- Set administrative boundaries
- Binds AS's together

Interior versus Exterior Routing Protocols

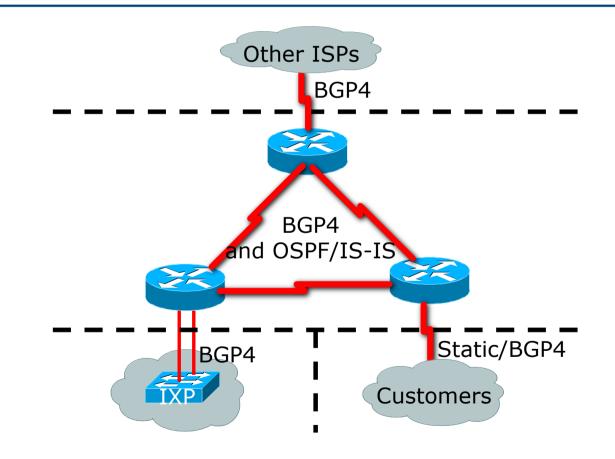
Interior

- Carries ISP infrastructure addresses only
- ISPs aim to keep the IGP small for efficiency and scalability

Exterior

- Carries customer prefixes
- Carries Internet prefixes
- EGPs are independent of ISP network topology

Hierarchy of Routing Protocols



FYI: Default Administrative Distances

Route Source	Cisco	Juniper	Huawei	Brocade	Nokia/ALU	Mikrotik
Connected Interface	0	0	0	0	0	0
Static Route	1	5	60	1	1	1
EIGRP Summary Route	5	N/A	?	N/A	N/A	N/A
External BGP	20	170	255	20	170	20
Internal EIGRP Route	90	N/A	?	N/A	N/A	N/A
IGRP	100	N/A	?	N/A	N/A	N/A
OSPF	110	10	10	110	10	110
IS-IS	115	18	15	115	18	N/A
RIP	120	100	100	120	100	120
EGP	140	N/A	N/A	N/A	N/A	N/A
External EIGRP	170	N/A	?	N/A	N/A	N/A
Internal BGP	200	170	255	200	130	200
Unknown	255	255	?	255	?	

Routing Basics

ISP Workshops